Password hashing

CS 161: Computer Security

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Passwords

Tension between usability and security

choose memorable passwords

choose råndom and long passwords (hard to guess)

Attack mechanisms

- Online guessing attacks
 - Attacker tries to login by trying different user passwords in the live system
- Social engineering and phishing
 - Attacker fools user into revealing password
- Eavesdropping
 - Network attacker intercepts plaintext password on the connection
- Client-side malware
 - Key-logger/malware captures password when inserted and sends to attacker
- Server compromise
 - Attacker compromises server, reads storage and learns passwords

Defences/mitigations

Network eavesdropper:

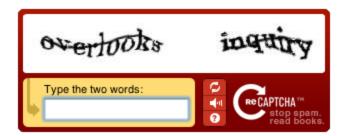
Encrypt traffic using TLS (will discuss later)

Client-side malware: hard to defend

- Intrusion detection mechanisms detect malware when it is being inserted into the network
- Various security software (e.g., anti-virus)
- Use two-factor authentication

Mitigations for online-guessing attacks

- Rate-limiting
 - Impose limit on number of passwords attempts
- CAPTCHAs: to prevent automated password guessing



 Password requirements: length, capital letters, characters, etc.

Mitigations for server compromise

- Suppose attacker steals the database at the server including all password information
- Storing passwords in plaintext makes them easy to steal
- Further problem: users reuse passwords at different sites!

Don't store passwords in plaintext at server!

Hashing passwords

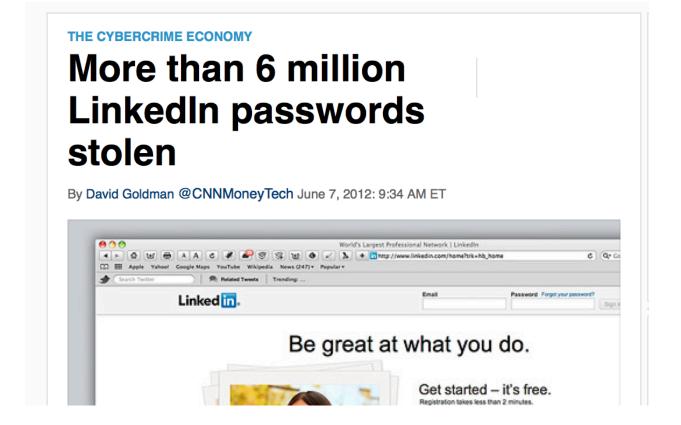
- Server stores hash(password) for each user using a cryptographic hash function
 - hash is a one-way function

username	hash of password
Alice	hash(Alice's password)
Bob	hash(Bob's password)

 When Alice logs in with password w (and provides w to server), server computes hash(w) and compares to Alice's record

Password hashing: problems

- Offline password guessing
 - Dictionary attack: attacker tries all passwords against each hash(w)
 - If D is dictionary size, n number of hashes passwords, attack takes Dn
 - Study shows that a dictionary of 2²⁰ passwords can guess 50% of passwords
- Amortized password hashing
 - Idea: One brute force scan for all/many hashes (D+n time)
 - Build table (H(password), password) for all 2²⁰
 passwords
 - Crack 50% of passwords in this one pass



LinkedIn was storing h(password)

"Link" was the number one hacked password, according to Rapid7. But many other LinkedIn users also picked passwords - "work" and "job" or example — that were associated with the career site's content.

Religion was also a popular password to pic — "god," "angel" and "jesus" also made the list.

Prevent amortized guessing attack

- Randomize hashes with salt
- Server stores (salt, hash(password, salt)), salt is random
- Two equal passwords have different hashes now
- Dictionary attack still possible, BUT need to do one brute force attack per hash now, not one brute force attack for many hashes at once
- Attacks takes Dn time instead of D+n time

Salted hash example

username	salt	hash of password
Alice	235545235	hash(Alice's password, 235545235)
Bob	678632523	hash(Bob's password, 678632523)

Attacker tries to guess Alice's password:

Computes table

```
'aaaaaa' hash('aaaaaaa', 235545235),
'aaaaab' hash('aaaaab', 235545235),
...
'zzzzzzz' hash('zzzzzzz', 235545235)
```

This table is useless for Bob's password because of different salt

Increase security further

- Would like to slow down attacker in doing a dictionary attack
- Use slow hashes = takes a while to compute the hash
- Define

```
H(x) = hash(hash(...hash(x))))
use with x = password || salt
```

- Tension: time for user to authenticate & login vs attacker time
- If H is 1000 times slower and attack takes a day with H, attack now takes 3 years with F

Conclusions

- Do not store passwords in cleartext
- Store them hashed with salts, slower hash functions better